

The purpose of this talk is to introduce *2-toposes*, which are 2-categories rich enough so that a lot of category theory can be developed inside them. They are one part of a meta-theory whose purpose is to organise the combinatorial machinery of higher dimensional category theory. In particular if you have a 2-topos \mathcal{K} then you can express what it means for an object of \mathcal{K} to have small colimits, and the basic theory of colimits unfolds in nice analogy with $\mathcal{K}=\mathbf{CAT}$ case.

Overview:

- operadic motivations.
- 2-toposes and examples.
- results on internal cocompleteness.

One can recognise a space X as having the homotopy type of a loop space when there are continuous functions

$$K_n \times X^n \rightarrow X$$

respecting certain axioms, where K_n is a convex polytope called the n -th associahedron.

The categorical/combinatorial ingredients for the “classical” theory of operads are:

1. a braided monoidal category \mathcal{V} (which in the above example is Top).
2. an operad in \mathcal{V} which is a sequence

$$K : \mathbb{N} \rightarrow \mathcal{V}$$

of objects of \mathcal{V} together with more data and axioms encoding substitution.

3. K can act on objects

$$X : 1 \rightarrow \mathcal{V}$$

of \mathcal{V} to express structure.

Colimits arise in the classical theory of operads in the following ways:

1. defining the substitution tensor product of collections.

$$(K \circ L)_n = \sum_{n_1 + \dots + n_r = n} K_r \otimes L_{n_1} \otimes \dots \otimes L_{n_r}$$

2. turning operads into monads.

$$\overline{K}(X) = \sum_{n \in \mathbb{N}} K_n \otimes X^{\otimes n}$$

3. comparing different types of operads (eg: symmetric versus non-symmetric).

all requiring that the \mathcal{V} you work in has enough colimits, and that these colimits are compatible with \otimes .

From the perspective of higher category theory the raw data for a space X is its underlying globular set

$$X_0 \xleftarrow{\quad} X_1 \xleftarrow{\quad} X_2 \xleftarrow{\quad} X_3 \xleftarrow{\quad} \dots$$

and the classical

$$K : \mathbb{N} \rightarrow \mathcal{V} \quad X : 1 \rightarrow \mathcal{V}$$

is replaced by

$$K : \text{Tr} \rightarrow \text{Sp}_{\mathbb{G}}(\text{Set}) \quad X : 1 \rightarrow \text{Sp}_{\mathbb{G}}(\text{Set})$$

Another important class of examples: take a \mathcal{V} from the classical theory and suspend it to get an n -globular category. **Organising the theory of higher operads takes you into the world of globular categories!**

A 2-cell

$$\begin{array}{ccc}
 A & \xrightarrow{g} & B \\
 & \searrow f & \nearrow h \\
 & C &
 \end{array}
 \quad \begin{array}{c}
 \phi \\
 \Rightarrow
 \end{array}$$

in a 2-category \mathcal{K} exhibits h as a *left extension* of f along g when $\forall k$ pasting with ϕ :

$$\kappa : h \Rightarrow k \quad \mapsto \quad \begin{array}{ccc}
 A & \xrightarrow{g} & B \\
 & \searrow f & \nearrow h \\
 & C &
 \end{array}
 \quad \begin{array}{c}
 \phi \\
 \Rightarrow
 \end{array}
 \quad \begin{array}{c}
 h \\
 \Downarrow \kappa \\
 k
 \end{array}$$

provides a bijection between 2-cells $h \Rightarrow k$ and 2-cells $f \Rightarrow kg$.

This 2-cell is a *pointwise* left extension when for all $b : X \rightarrow B$, the composite

$$\begin{array}{ccc}
 g/b & \longrightarrow & X \\
 \downarrow & \Rightarrow & \downarrow b \\
 A & \xrightarrow{g} & B \\
 \searrow f & \Downarrow \phi & \swarrow h \\
 & C &
 \end{array}$$

exhibits hb as a left extension.

Roughly then, C is cocomplete when pointwise left extensions into it exist.

So we can ask:

1. Is $\mathrm{Sp}_{\mathbb{G}}(\mathrm{Set})$ or more generally $\mathrm{Sp}_{\mathbb{G}}(\mathcal{V})$ where \mathcal{V} is a cocomplete category, cocomplete as a globular category?
2. What about $\Sigma^n(\mathcal{V})$?
3. Is $\mathrm{Sp}_{\mathbb{G}}(\mathrm{Set})$ in fact the small globular colimit completion of 1 just as Set is the small colimit completion of 1 ?

Recall that a category \mathcal{E} is an *elementary topos* when

1. it has finite limits,
2. is cartesian closed, and
3. has a subobject classifier.

So what's a 2-categorical analogue of subobject classifier?

To any functor $f : A \rightarrow \text{Set}$ a version of the Grothendieck construction associates to f a discrete opfibration $G(f) : e(f) \rightarrow A$ to provide a fully faithful functor

$$G_{U,A} : \text{CAT}(A, \text{Set}) \rightarrow \text{DOPFib}(A)$$

whose image consists of those discrete opfibrations with small fibres.

The analogy with topos theory comes about because G can be described as

$$\begin{array}{ccc}
 e(f) & \longrightarrow & \text{Set} \bullet \\
 \downarrow G(f) & \lrcorner & \downarrow U \\
 A & \xrightarrow{f} & \text{Set}
 \end{array}$$

Definition 1 *Let \mathcal{K} be a finitely complete 2-category. A discrete opfibration $\tau : \Omega_{\bullet} \rightarrow \Omega$ in \mathcal{K} is classifying when the functors*

$$G_{\tau, A} : \mathcal{K}(A, \Omega) \rightarrow \mathbf{DFib}(1, A)$$

are fully faithful for all $A \in \mathcal{K}$.

Definition 2 *A 2-topos $(\mathcal{K}, (-)^{\circ}, \tau)$ is a finitely complete cartesian closed 2-category \mathcal{K} equipped with a duality involution $(-)^{\circ}$ and a classifying discrete opfibration $\tau : \Omega_{\bullet} \rightarrow \Omega$.*

Theorem 3 *Let*

$$\mathcal{A} \begin{array}{c} \xleftarrow{E} \\ \perp \\ \xrightarrow{S} \end{array} \mathcal{B}$$

be a 2-adjunction and \mathcal{A} and \mathcal{B} be finitely complete 2-categories. If

- 1. E preserves pullbacks, and*
- 2. naturality squares for the counit at maps in \mathcal{A} which are discrete opfibrations, are pullbacks*

then S preserves classifying discrete opfibrations.

Examples:

1. The subobject classifier

$$\tau : 1 \rightarrow \Omega$$

of an elementary topos \mathcal{E} regarded as an internal functor is a classifying discrete opfibration in $\text{Cat}(\mathcal{E})$.

2. $U : \text{Set}_\bullet \rightarrow \text{Set}$ is a c.d.o for CAT . Similarly with Set replaced by finite sets.
3. $\text{Sp}_{\mathbb{G}}(U)$, $\Sigma \text{Sp}_{\mathbb{G}}(U)$, $\Sigma^2 \text{Sp}_{\mathbb{G}}(U)$, ... are c.d.o's for $\text{CAT}(\hat{\mathbb{G}})$.
4. there's a version of $\text{Sp}_{\mathbb{G}}(U)$ for $\text{CAT}(\hat{\mathbb{C}})$.

Definition 4 *A map $f : A \rightarrow B$ in a 2-topos is admissible when the discrete fibration f/B is in the image of the canonical functor*

$$\mathcal{K}(B, \hat{A}) \rightarrow \text{DFib}(A, B)$$

An object A is admissible when the map 1_A is admissible. A is small when it is admissible and \hat{A} is admissible.

There is a general result which enables one to characterise smalls and admissibles. For instance for $\text{Sp}_{\mathbb{G}}(U)$ the corresponding small globular categories are just those that are small at each level.

Definition 5 *An object A in a 2-topos is co-complete when it admits all pointwise left extensions along any $h : X \rightarrow Z$ such that X is small and Z is admissible.*

Theorem 6 *Let*

$$\mathcal{A} \begin{array}{c} \xleftarrow{E} \\ \perp \\ \xrightarrow{S} \end{array} \mathcal{B}$$

be a 2-adjunction, \mathcal{A} and \mathcal{B} be 2-toposes which exhibit small lax pullbacks. If

- 1. E preserves small objects, admissible objects and lax pullbacks up to a right adjoint section; and*
- 2. S preserves admissible objects*

then S preserves objects which are both admissible and cocomplete.

Theorem 7 1. $\text{Sp}_{\mathbb{G}}(\text{Set})$ is the small globular colimit completion of $\mathbb{1}$.

2. $\text{Sp}_{\mathbb{G}}(\text{Set})$ is cartesian closed as a globular category.

In fact the above results are true when \mathbb{G} is replaced by an arbitrary small category \mathbb{C} .